cgroup2 Resource Control strategies w/ resctl-demo

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The Goal of resource control

Work-conserving full-OS resource isolation

What does that mean?

- Work-conserving:
 - Don't keep machine idle if there's work to do
- Full-OS:
 - Transparent
 - Keep doing what you've been doing and overlay isolation
 - No need for direct IO, hard-allocate mem, separate FS







	cpu%	mem	swap	rbps	wbps	cpuP%	memP%	ioP%
workload	92.3	28.7G	856K	15.0M	3.9K	84.2	-	0.0
sideload	-	-	_	_	_	-	-	-
hostcritical	7.4	647M	_	-	_	31.7	-	0.0
system	0.0	25.8M	1.5M	-	-	0.1	-	-
user	0.0	577M	3.4M	-	-	0.0	-	-
	100	30.6G	806M	15.0M	3.9K	86.7	_	0.0

[comp.cgroup] Cgroup and Resource Protection - 'i': index, 'b': back

Cgroup and Resource Protection

<u>Scenario</u>: A web server, an external memory leak, and no resource control

Imagine a fleet of web servers whose main job is running a web application to service user requests. But like all systems, a lot of other apps need to run: The system software and the web app require regular updates; fleet-wide maintenance tools keep the configuration in compliance; various monitoring and alarm frameworks need to run - and so on.

Let's say a maintenance program started by cron is malfunctioning and keeps leaking memory, and I that the malfunction is wall-clock dependent: This is a bad scenario that would stay hidden I during testing, and then trigger at the same time in production. What would happen to the I system and fleet without resource control enabled?

We can simulate this scenario with the rd-hashd workload simulator, and a test program intentionally designed to leak memory. rd-hashd is already started targeting the full load. Give it some time to ramp up, and let its memory footprint grow to closely fill up the machine. Check out the memory utilization graph by pressing 'g'.

Once the workload sufficiently ramps up, you can see how the system behaves without any protection, by selecting the button below to disable all resource control mechanisms and start a memory-hog program. The problematic program will start as rd-sysload-memory-hog.service under system:

[Disable all resource control features and start memory hog]

MARVING: Because the system is running without any protection, nothing can guarantee the system's responsiveness. Everything, including this demo program, will get sluggish, and might completely stall. Depending on how the kernel OOM killer reacts, the system may or may not recover in a reasonable amount of time. To avoid the need for a reboot, stop the experiment once the system becomes sluggish:

[Stop the memory hog and restore resource control]

That wasn't ideal, was it? If this were a production environment and the failure happened across a large number of machines at the same time, our users would notice the disturbances, and the whole site might go down.



The components

- rd-hashd: Latency sensitive main workload
- Sysloads: Possibly misbehaving secondary workloads
- Sideloads: Secondary workloads running under sideloader
- rd-agent: System and experiment management
- resctl-demo: TUI for doc, monitoring and contol

[intro.sysreqs] System Re
<pre>* SysReq::IoCost: blk-iocost is the comprehensively control IO capaci with CONFIG_BLK_CGROUP_IOCOST. Fo Articles/793460/</pre>
<pre>* SysReq::IoCostVer: blk-iocost reconstrol quality and visibility du with these updates is recommended Articles/830397/</pre>
<pre>* SysReq::NoOtherIoControllers: Othe - can interfere and shouldn't have</pre>
If configured through systemd, rep and IoDeviceLatencyTargetSec conf
<pre>* SysReq::AnonBalance: Kernel memory during the v5.8 development cycle footing with page cache and made details: https://lwn.net/Articles</pre>
* SysReq::Btrfs: Working IO isolatic avoid priority inversions. Current filesystem.
The OS must be installed with btr
* SysReq::BtrfsAsyncDiscard: Many St discards are issued in bulk, which Async discard is a btrfs feature discards.
It can be enabled with "discard=a available, resctl-demo will autom mount option. For details: https:
<pre>* SysReq::NoCompositeStorage: Curre and md, break the chain of custod control and cause severe priority</pre>
The filesystem must be on a physi
* SvsRea::IoSched: bfa IO scheduler

equirements – 'i': index, 'b': back
e new IO controller which can ity distribution proportionally. Enabled or details: https://lwn.net/
<pre>ceived significant updates to improve uring the v5.10 development cycle. A kernel d. For details: https://lwn.net/</pre>
her IO controllers – io.max and io.latency ve active configurations.
emove all IO{Read Write}{Bandwidth IOPS}Max figurations.
ry management received a major update e which put anonymous memory on an equal swap useful, especially on SSDs. For s/821105/
ion requires support from filesystem to ntly, btrfs is the only supported
rfs as the root filesystem.
SSDs show significant latency spikes when ch can lead to severe priority inversions. that paces and reduces the total amount of
async" mount option on kernels >= v5.6. If matically remount the filesystem with the ://lwn.net/Articles/805300/
ently, composite block devices, such as dm dy for IOs, allowing cgroups to escape IO y inversions.
ical device.
r's implementation of proportional IO



Protection

Scenario

- Web server is running full-tilt
- A management job running from chef has a new bug leaking memory.
- Can we survive if the same fleet at the same time?

tilt g from chef has a new bug

Can we survive if the same behavior happens across the



workload	•						memP% -	
sideload	_	_	—	_	_	_	_	—
hostcritical	7.4	744M	_	_	19.7K	34.0	_	0.1
system	0.0	17 . 7M	1 . 5M	_	_	0.0	_	—
user	0.0	517M	3.4M	_	_	0.0	_	_
-	100	27 . 3G	678M	34 . 0M	23 . 7K	93.0	-	0.0

[comp.cgroup] Cgroup and Resource Protection – 'i': index, 'b': back –

Cgroup and Resource Protection

Scenario: A web server, an external memory leak, and no resource control

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[Stop the memory hog and restore resource control]

That wasn't ideal, was it? If this were a production environment and the failure happened across a large number of machines at the same time, our users would notice the disturbances, and the whole site might go down.





workload							memP% 66.3	
sideload	_	_	_	_	_	_	_	—
hostcritical	5.7	707M	44.0K	_	_	1.3	_	0.0
system	0.9	5.7G	1.0G	_	34.2M	0.0	_	—
user		85.6M	3 . 1M	_	_	0.0	_	_
-	14.2	31 . 0G	1 . 8G	63 . 7M	227M	1.4	63.2	60.3

[comp.cgroup] Cgroup and Resource Protection - 'i': index, 'b': back -

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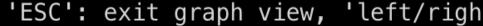
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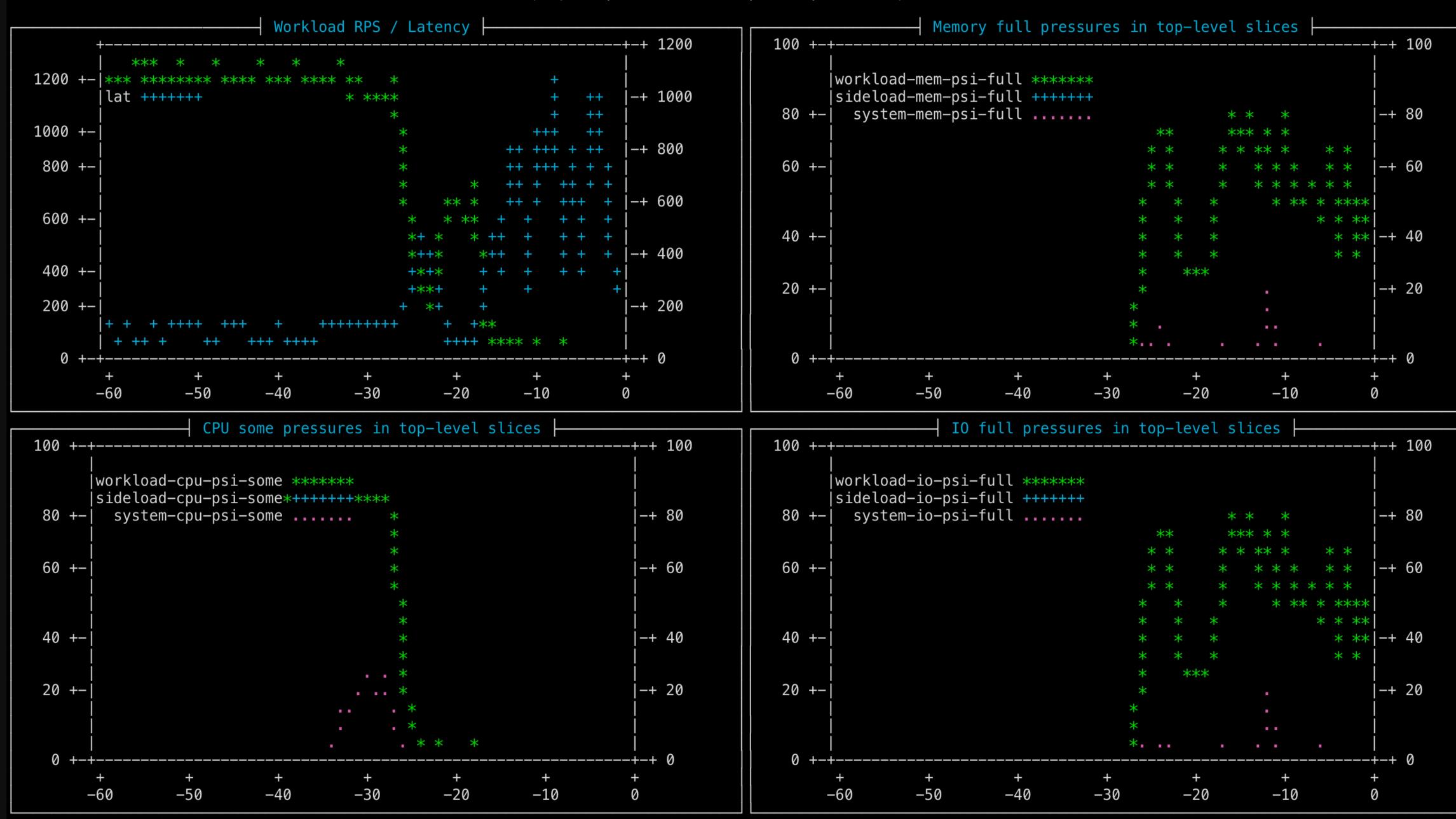
[Stop the memory hog and restore resource control]

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[rps/psi] | utilization | I0 | iocost/psi-some

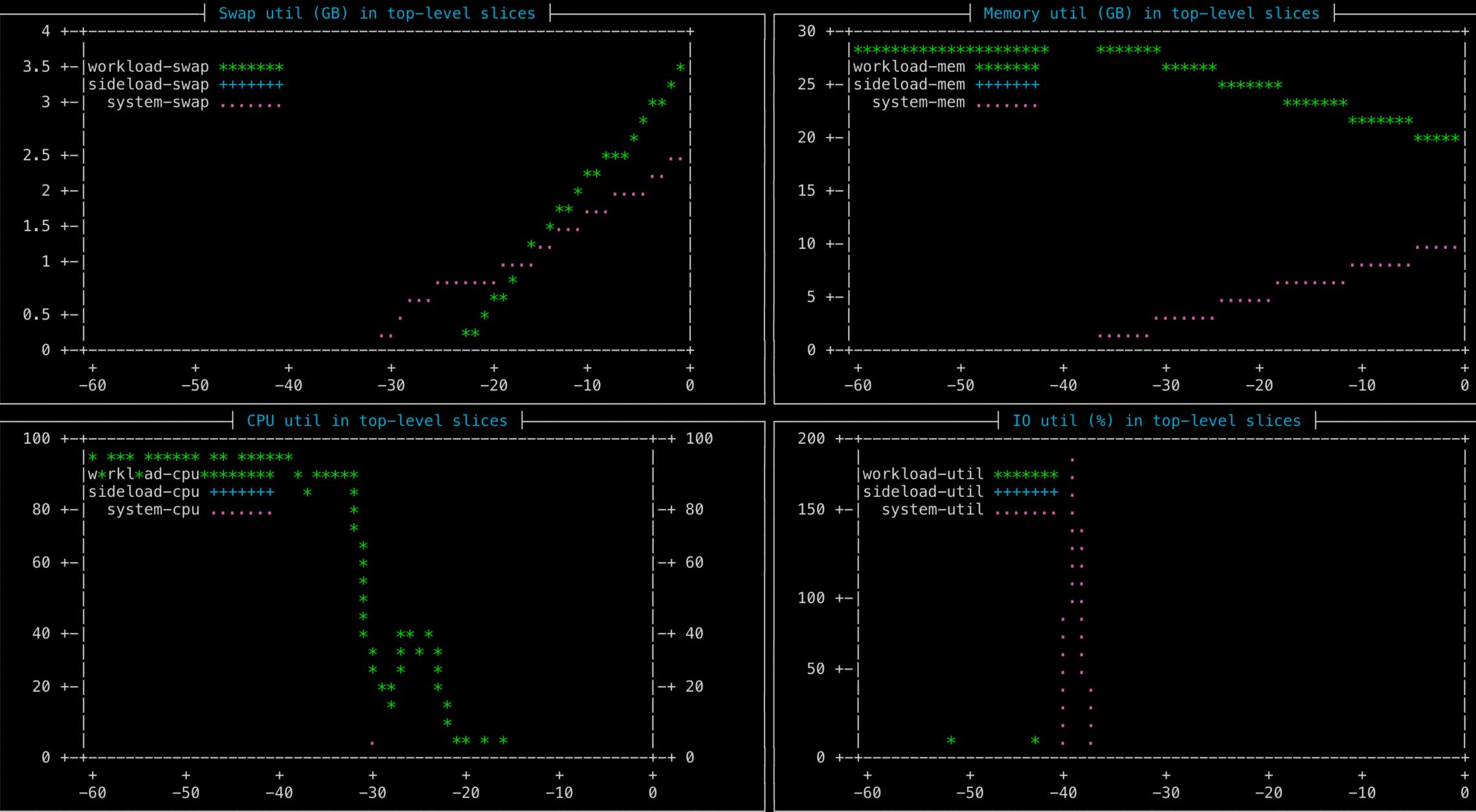


'ESC': exit graph view, 'left/right': navigate tabs, 't/T': change timescale



'ESC': exit graph view, 'left/right': navigate tabs, 't/T': change timescale

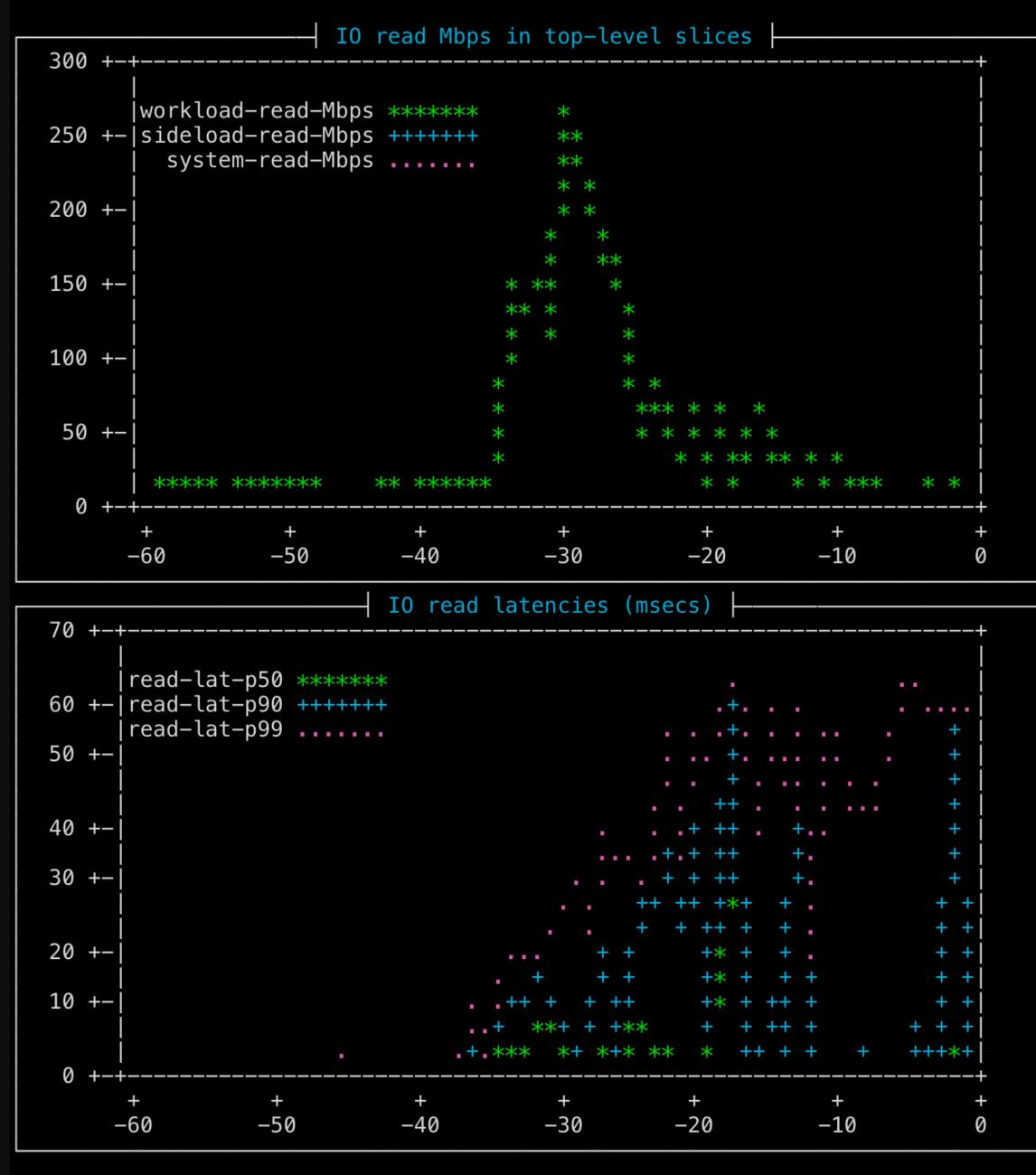
rps/psi [utilization] I0 iocost/psi-some





'ESC': exit graph view, 'left/right': navigate tabs, 't/T': change timescale

rps/psi | utilization | [I0] | iocost/psi-some



______ IO read Mbps in top-level slices _______ 250 +-+------ IO write Mbps in top-level slices ______ workload-write-Mbps ****** |sideload-write-Mbps +++++++ * * 200 +-| system-write-Mbps ** * * * * *** **** * ** * * * * *** * . 150 +-. . . * *** * ** * * ** * * 100 +-| 50 +-| * * . . ** . * Ø +-+-----++ + +-50 -40 -60 -30 -20 -100 IO write latencies (msecs) 80 +-+------70 +-|write-lat-p50 ****** . . |write-lat-p90 +++++++ . + . . 60 +-|write-lat-p99 . ++... . .. ++ + +*** +* + + + ++ + 50 +-** *+* * * * + + * * * * * * 40 +-+ + * * * + * 30 +-+ +* + +++ +** 20 +-. . + . . + ** + * ** 10 +-* *** * *** * * * * 0 +-+--+++++ +-50 -40 -30 -20 -100 -60



Same scenario with protection cpu.weight, memory.low and io.weight configured

cpu.weight, memory.low ar oomd on watch



workload	•						memP% 2.8	
sideload	—	_	_	—	_	—	_	_
hostcritical	6.6	721M	_	_	_	18.7	_	0.0
system	0.0	1.6G	345M	10.6K	38.0M	0.3	84.1	84.1
user	0.0	84.8M	21.2M	_	_	0.0	_	_
—	97.1	31 . 0G	7.2G	111M	48 . 1M	67.7	4.2	2.0

[comp.cgroup] Cgroup and Resource Protection - 'i': index, 'b': back slices.

You can check out what cgroups are on your system by running `systemd-cgls`. The kernel interface is a pseudo filesystem mounted under /sys/fs/cgroup.

<u>Turning on resource protection</u>

Once the kernel understands what applications are on the system and how they're grouped in cgroups, it can monitor resource consumption and control distribution along that hierarchy.

We'll repeat the same test run we did above, but leaving all the resource protection configurations turned on.

rd-hashd should already be running at full tilt. Once <u>workload</u>'s memory usage stops growing, let's start the same memory hog, but without touching anything else:

[Start memory hog]

Monitor the RPS and other resource consumption in the graph view ('g'). The RPS may go down a little and dip occasionally, but it'll stay close to full capacity no matter how long you let the memory hog go on. Eventually the memory hog will be killed off by OOMD. Try it multiple times by clearing the memory hog with the following button, and restarting it with the start button above:

[Stop memory hog]

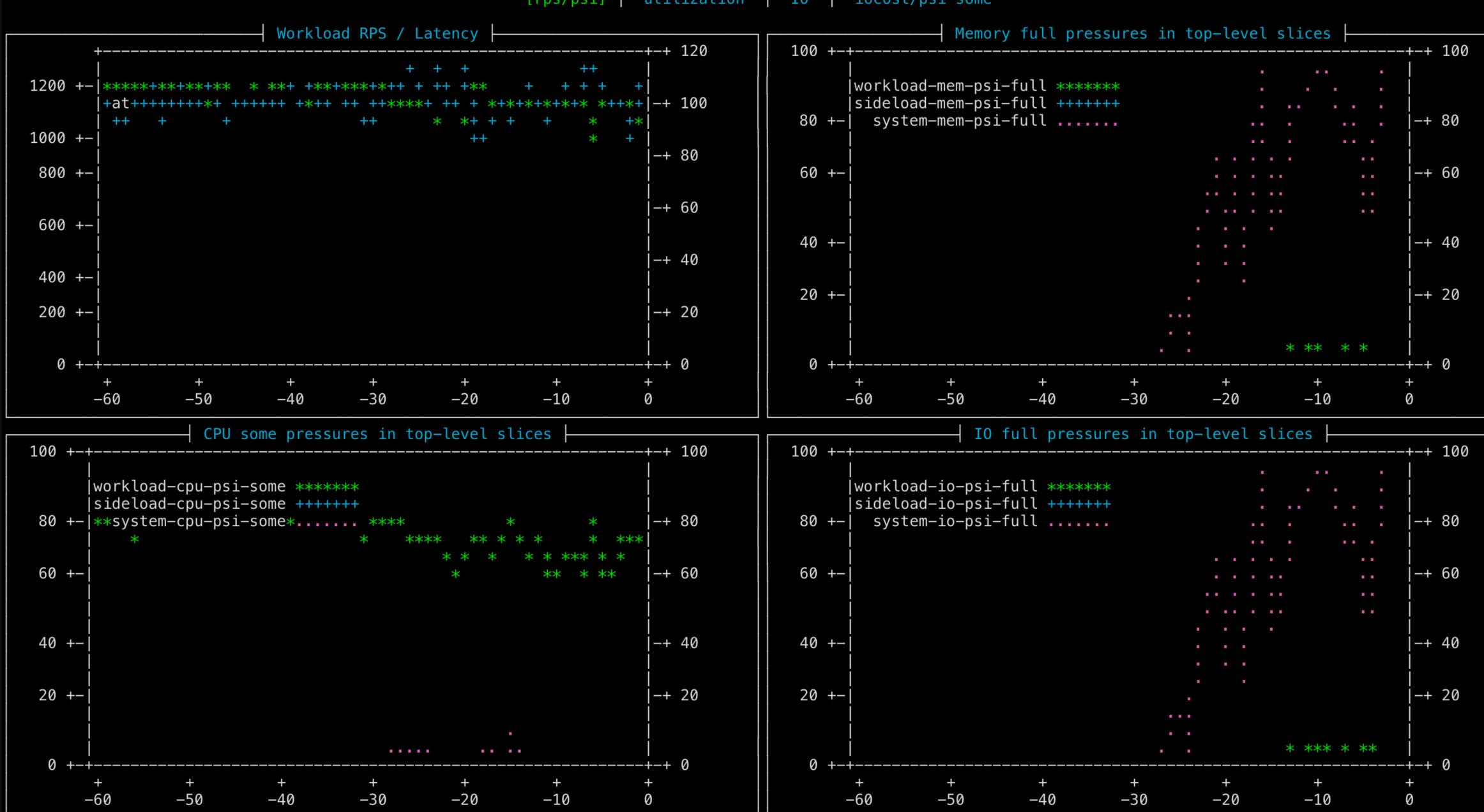
In this scenario, with resource control on, our site's not going down, and users probably aren't even noticing. If we have a working monitoring mechanism, the OOMD kills will raise an alarm, people can quickly find out which program was malfunctioning, and hunt down the bug. What could have been a site-wide outage got de-escalated into an internal oops.

This is cgroup resource protection at work. The kernel understood which part was important and protected its resources at the expense of the memory hog running in the low-priority portion of the system. Eventually, the situation for the memory hog became unsustainable and OOMD terminated it.

We'll go into details in the following pages but here's a brief summary of how resource protection is configured in this demo.

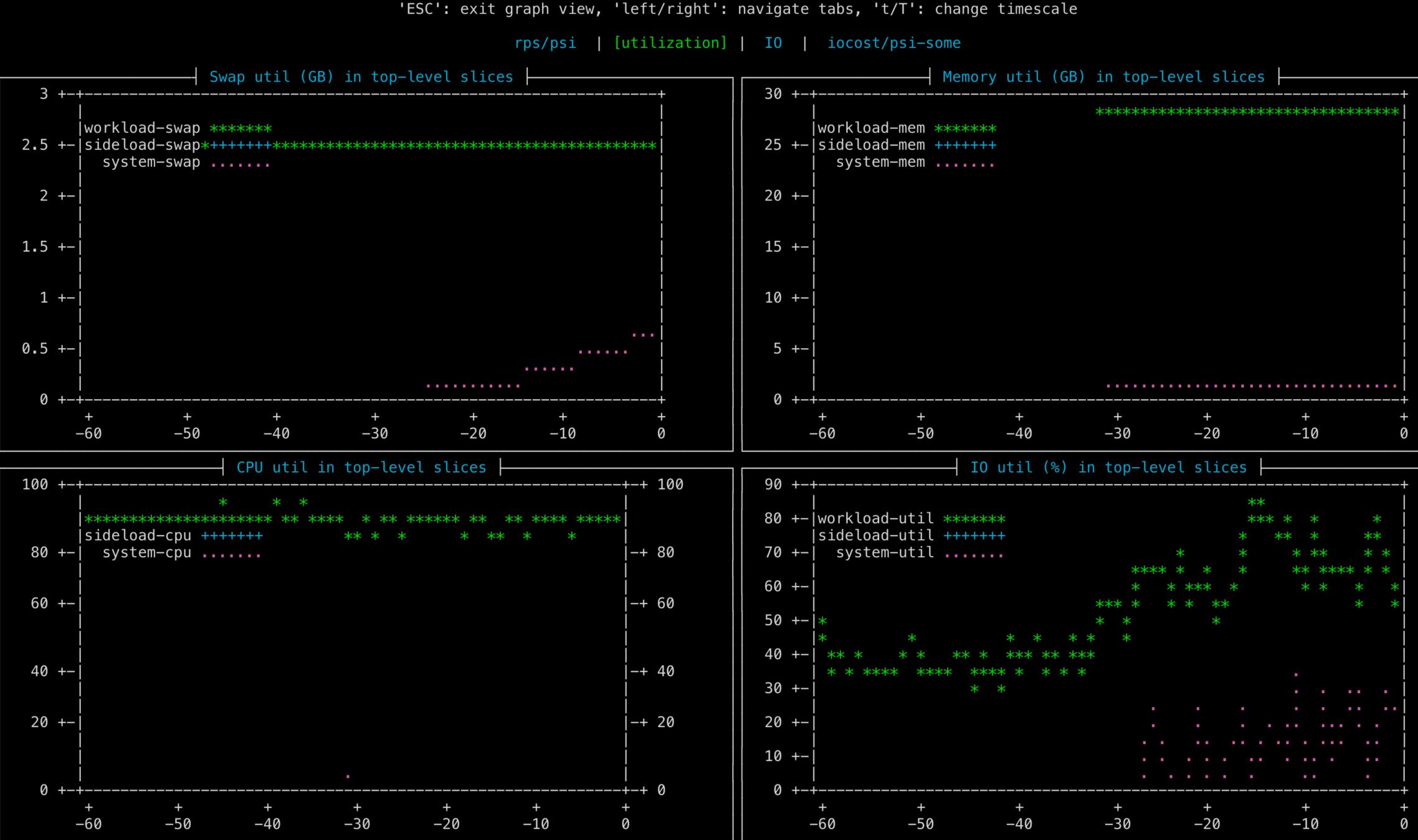


'ESC': exit graph view, 'left/right': navigate tabs, 't/T': change timescale



[rps/psi] | utilization | I0 | iocost/psi-some

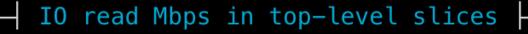


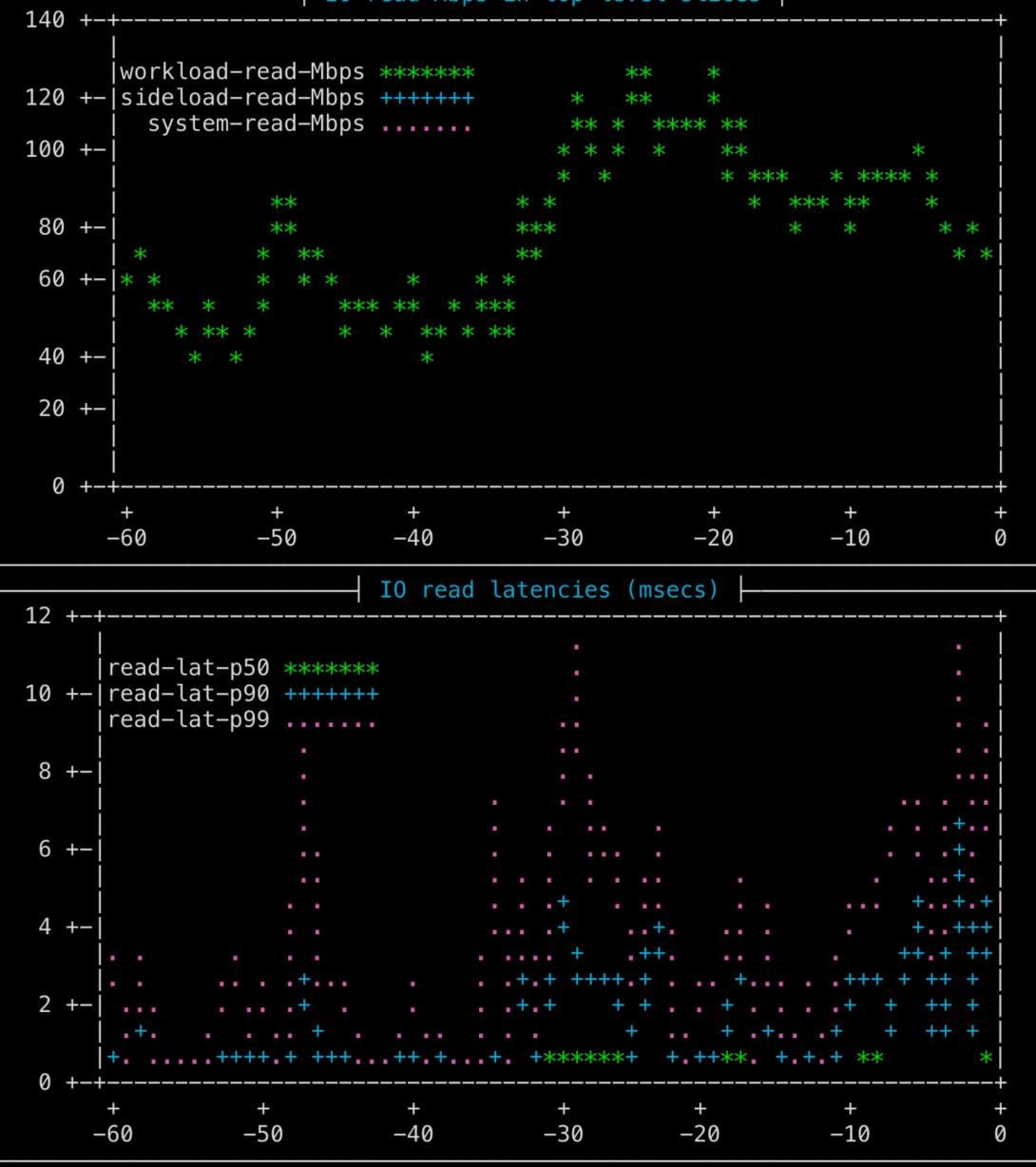




'ESC': exit graph view, 'left/right': navigate tabs, 't/T': change timescale

rps/psi | utilization | [I0] | iocost/psi-some

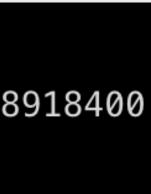




I0 write Mbps in top-level slices 45 +-|workload-write-Mbps ****** . * sideload-write-Mbps +++++++ . 40 +-| system-write-Mbps 35 +-1 30 +-1 * ** 25 +-| * * * * ** 20 +-* * * * ** ** 15 +- | ** * * * * *** ** * * * * * * * * * * 10 +-| *** *** ** * * * * * * * * * * * ** *** ** * * ** * * * * * * * 5 +-|* ** * * * Ø +-+------+ -20 -50 -40-30 -10-60 0 IO write latencies (msecs) 14 +-+----write-lat-p50 ****** 12 +-|write-lat-p90 +++++++ |write-lat-p99 10 +-8 +-1 . . 6 +-|. .+. 4 +- | -.+.+.+ +++ + .+. . . .++. ...+. .. . 2 +- | . . + . .+..+. .+. +++ *+ + + + + ++++ + *+* 0 + - + - --60 -50 -30 -20 -100 -40



[Jan 22 16:34:53 rd-oomd] [../src/oomd/plugins/BaseKillPlugin.cpp:319] Killed 1: 112038(memory-growth.p)
[Jan 22 16:34:53 rd-oomd] [../src/oomd/plugins/BaseKillPlugin.cpp:319] Killed 1: 112038(memory-growth.p) [Jan 22 16:34:54 rd-oomd] [../src/oomd/Log.cpp:117] 36.84 22.28 6.04 system.slice/rd-sysload-memory-hog.service 1058918400 ruleset:[protection against heavy system.slice thrashing] detectorgroup:[Sustained thrashing in system.slice] killer:kill_by_memory_size_or_growth v2



Sideloading

What's sideloading?

- Machines can't be utilized fully for reasons such as scheduling inefficiency and DR buffer
- Maybe simply a generalized case of protection?





workload	•						memP% -	
sideload	—	—	—	_	—	—	—	—
hostcritical	6.4	741M	32 . 0K	_	_	10.6	_	0.0
system	43.4	2 . 9G	792K	32.0K	143K	50.5	_	0.0
user	0.1	84 . 9M	21 . 0M	_	_	0.2	—	—
—	100	30.0G	578M	1.6M	5 . 5M	67.0	—	0.0

[side.intro] What is Sideloading? - 'i': index, 'b': back

<u>A naive approach</u>

In the previous chapter, we demonstrated that resource control can protect the main workload from the rest of the system. While rd-hashd was running at full load, we could throw many types of misbehaving workloads at the system with only limited impact on the main workload. We also saw that CPU control couldn't protect rd-hashd's latency from a CPU hog when rd-hashd was running close to full load.

If the CPU control can protect the main workload at moderate load levels, maybe we can simply run the sideloads under the same resource control configuration and stop them when the main workload's load spikes. Let's see how latency at 60% load and ramping up from there are impacted by running sideloads this way.

rd-hashd should already be running at 60% load. Once it's warmed up, start a Linux build job with 1x CPU count concurrency. Depending on the hardware and kernel configurations, rd-hashd may fail to hold the RPS with the usual 100ms latency target, so let's relax the latency target to 200ms too. Pay attention to how the latency in the left graph pane changes:

[Relax rd-hashd latency target and start linux build job]

Note how RPS is holding but latency deteriorates sharply. Press 'g' and check out the resource pressure graphs. Even though this page set the CPU weight of the build job only at a hundredth of rd-hashd, CPU pressure is noticeable. We'll get back to this later.

Now let's push the load up to 100% and see whether its ability to ramp up is impacted too:

[Set full load]

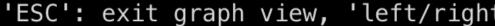
It climbs, but seems kind of sluggish. Let's compare it with a load rising but without the build job:

[Stop linux build, reset latency target and set 60% load]

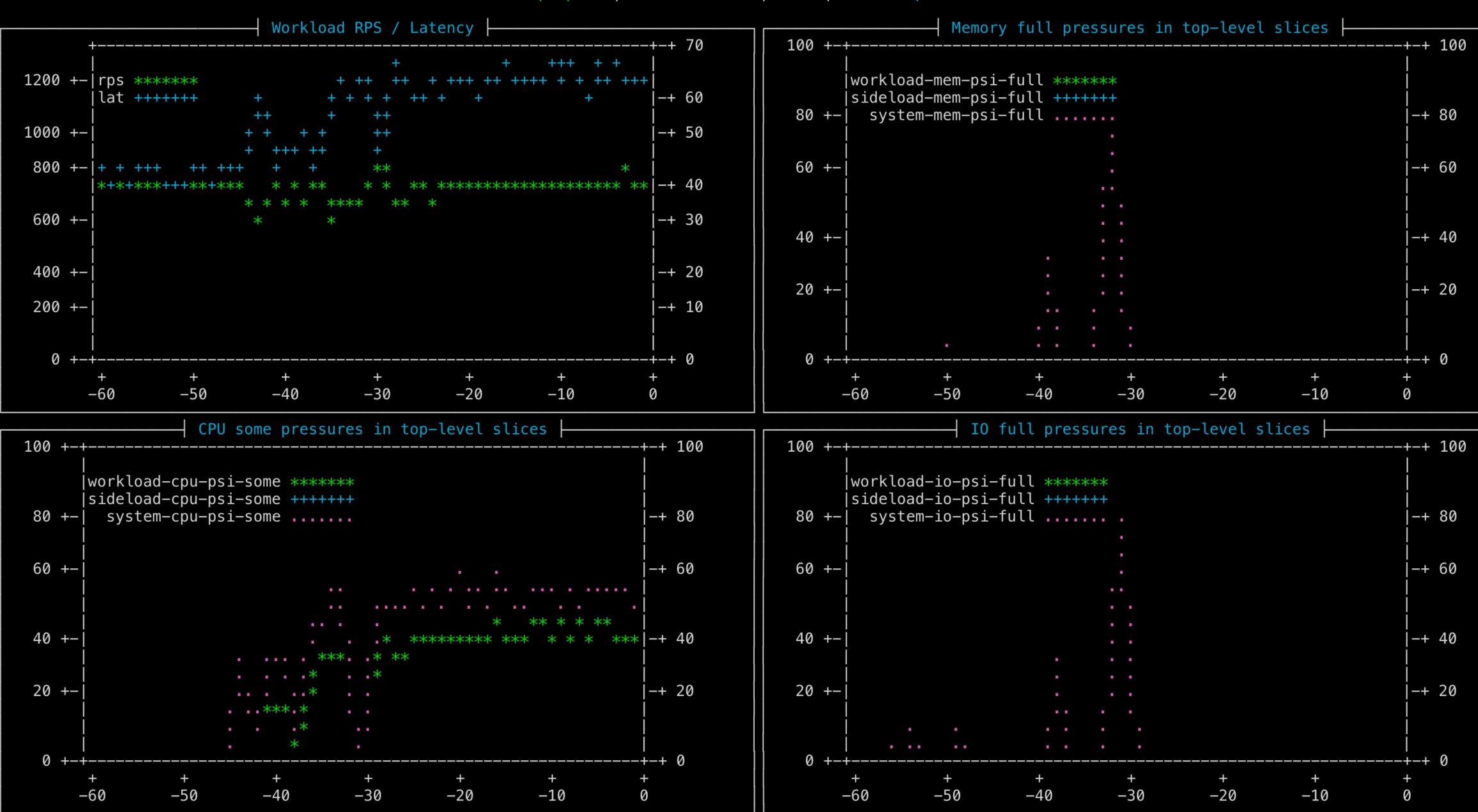
Wait until it stabilizes, then ramp it up to 100%:

[Set full load]





[rps/psi] | utilization | I0 | iocost/psi-some

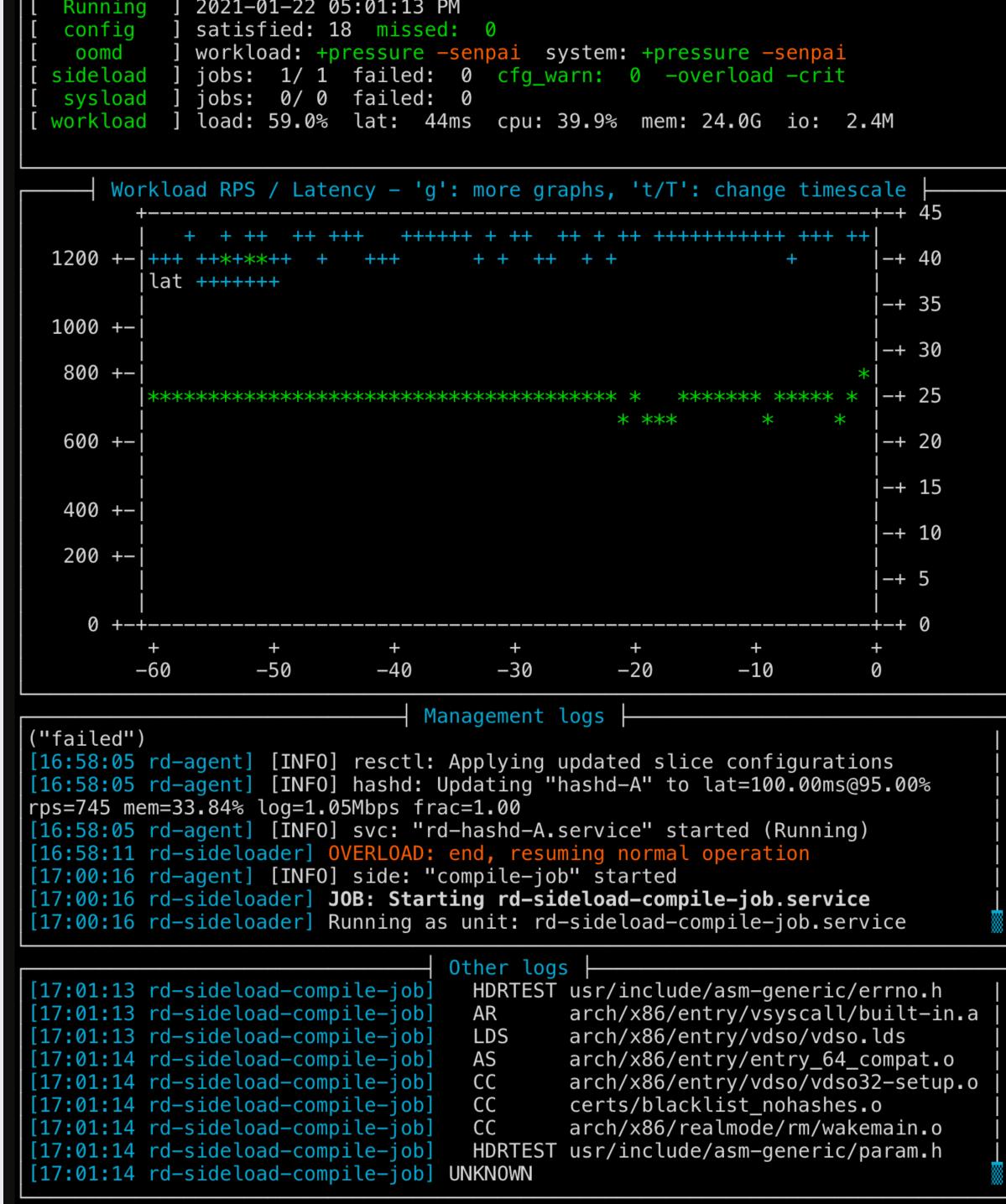


'ESC': exit graph view, 'left/right': navigate tabs, 't/T': change timescale



Why the latency deterioration?

- Scheduling inefficiencies
- CPU sub-resource contention. The same instructions take longer to execute as CPU gets saturated.
- Sideloader regulates cpu.max to maintain sufficient CPU headroom.



	cpu%	mem	swap	rbps	wbps	cpuP%	memP%	ioP%
workload	40.4	24.9G	772K	2.7M	4.3M	5.9	—	0.1
sideload	13.6	3 . 0G	_	32 . 0K	32 . 0M	39.3	34.7	34.7
hostcritical							3.9	3.9
system	—	174M	788K	_	—	—	—	—
user	0.0	84 . 9M	21 . 0M	_	—	0.0	_	_
—	64.5	29 . 5G	578M	2 . 7M	36.3M	43.9	16.1	16.3

[side.sideloader] The Sideloader – 'i': index, 'b': back –

Now, let's repeat the experiment from the last section, but launch the Linux compile job as a sideload that's supervised by the sideloader.

rd-hashd should already be running at 60% load. Once it warms up, let's start a Linux compile job with 1x CPU count concurrency as before. It'll have the same resource weights as before; the only difference is that it's now running under the supervision of sideloader:

[Start linux compile job as a sideload]

While the Linux source tree is being untarred, depending on the memory situation and IO performance, you may see a brief spike in latency as the kernel tries to figure out the hot working set of the primary workload. Once the compile jobs start running, <u>sideload</u> starts consuming CPU time and pushing up CPU utilization. Look in the upper right pane, and at the utilization graphs in the graph view ('g'). CPU utilization will go up but won't reach 100%.

Look at the latency graph. It's gone up a bit, but a lot less than before, when we ran it as a sysload. You can change the time scale with 't/T' to compare the latencies before and after. How much sideloads impact latency is determined by the headroom: The bigger the headroom, the lower the CPU utilization, and the lower the latency impact.

The current state, in which the sideload fills up the machine's left-over capacity, with a controlled latency impact on the main workload – can be sustained indefinitely. Let's see how ramping up to full load behaves:

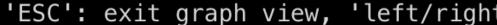
[Set full load]

The difference from the no sideload case should be significantly less pronounced. Look at how the sideloader state transitions to overload, freezing the Linux compile job. The build job is configured to expire after being frozen for 30s. Let's wait and see what happens.

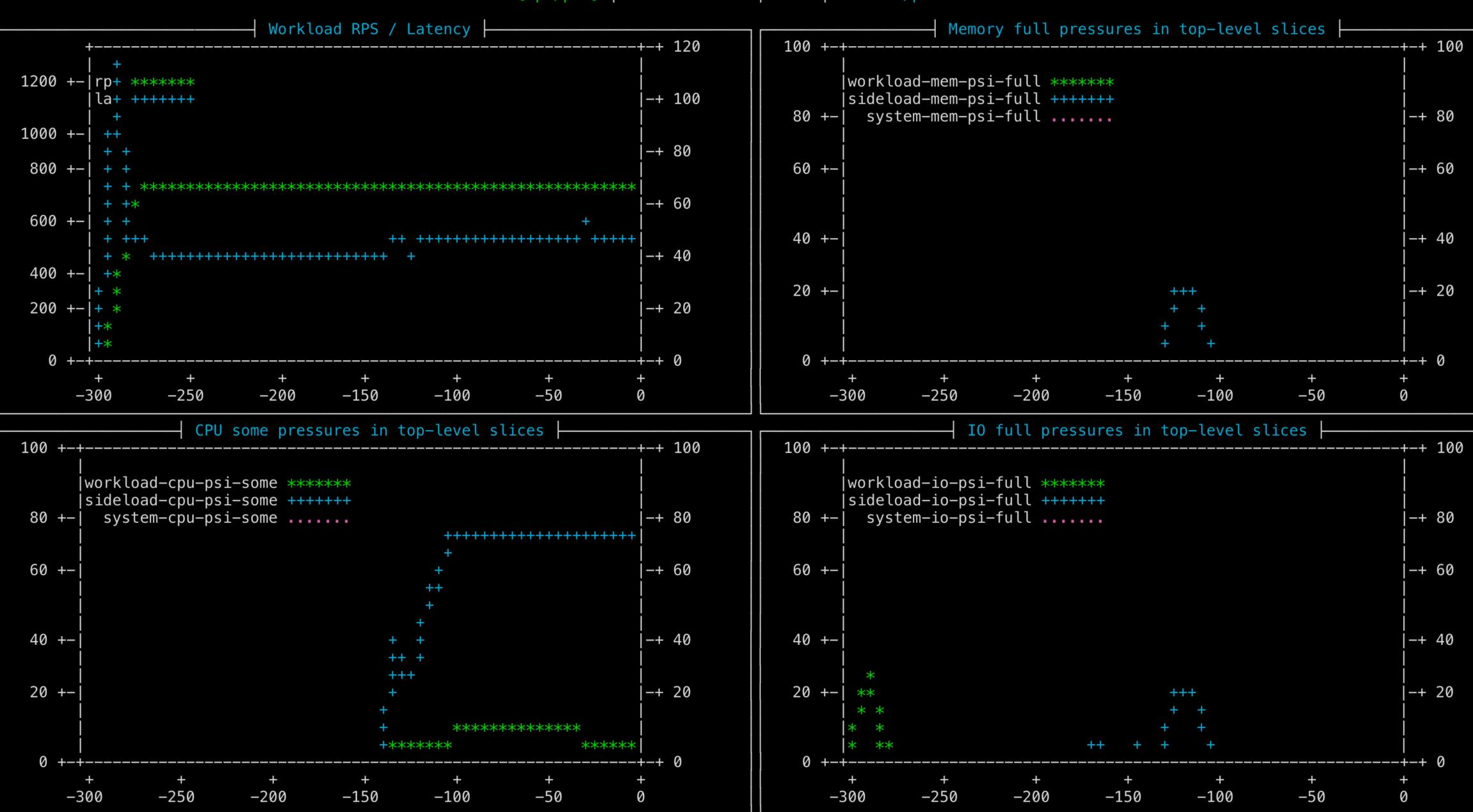
We just demonstrated that, with the help of sideloader, sideloads can utilize the left-over capacity of the machine, with only a controlled and configurable latency impact on the main workload, and without significantly impacting the main workload's ability to ramp up when needed.

<u>Read on</u>

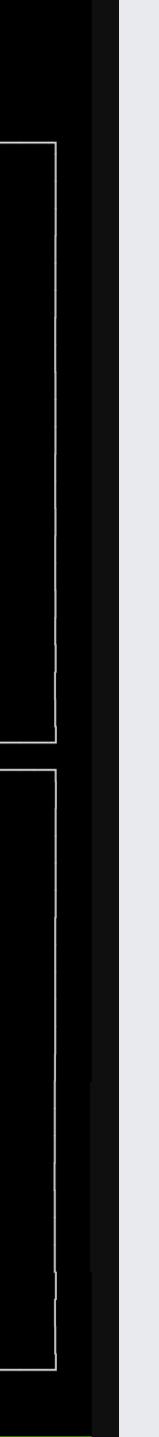


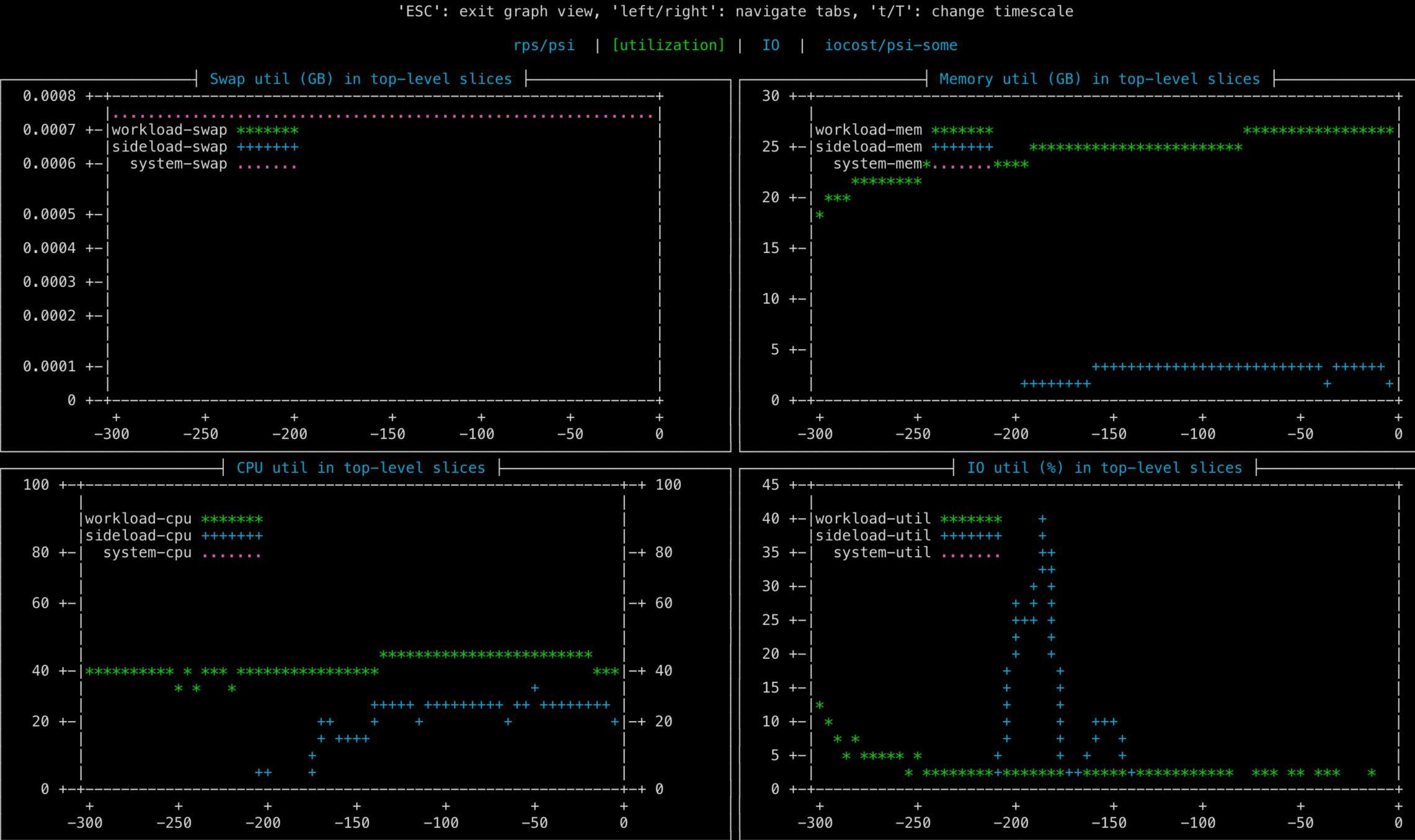


[rps/psi] | utilization | I0 | iocost/psi-some



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Under active development

- Matching resource behavior of production workload
- resctl-bench: Whole system benchmarks for tuning,
 - resource control evaluations and verifications

or of production workload n benchmarks for tuning, ns and verifications

https://facebookmicrosites.github.io/resctl-demo-website



